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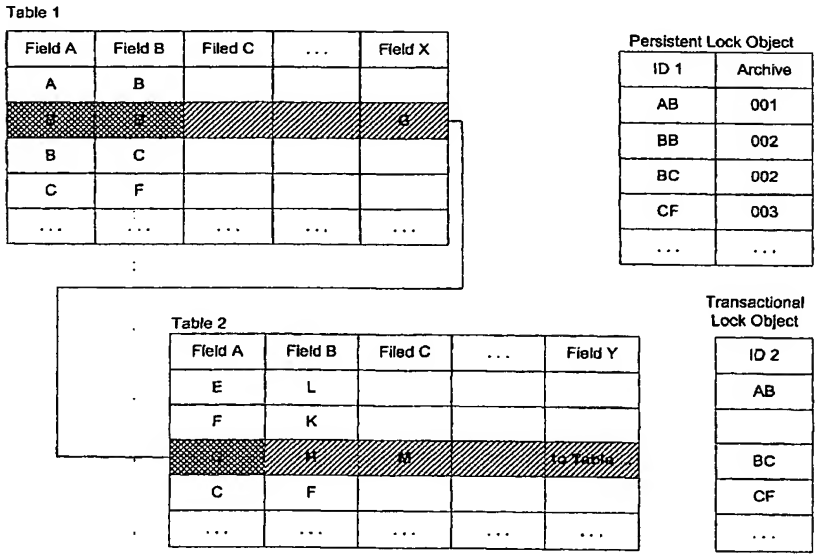
(54) Title: METHODS AND SYSTEMS FOR MOVING DATA USING LOCKS



(57) Abstract: The Invention relates to a process for moving data objects in a computer system from a first to a second storage location, comprising: a) selecting one or more data objects having an identifier (ID) from the first storage location, b) storing said ID in a first lock object, c) storing said ID in a second lock object, d) storing a data object, the ID of which is contained in the first lock object, at the second storage location, e) deleting a data object, the ID of which is contained in the first lock object, from said first storage location, f) deleting an ID from the first lock object earliest at a time at which step c) for the respective data object assigned to that ID has been completed, g) deleting an ID from the second lock object earliest at a time at which step b) for a particular ID has been completed.

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Methods and Systems for Moving Data Using Locks

Background of the Invention

Field of the Invention.

The technical field of this invention is in the area of electronic data processing. More particularly, the invention relates to methods, computer program products and systems for data moving.

Description of the Related Art

Moving of data objects is well known to every user of a computer and is a standard procedure, which is routinely applied. A special application of moving data objects is the archiving process, by which data objects are moved from a first to a second storage location for safety and/or performance reasons. In enterprises, enterprise resource planning software (ERP) applications are used to control or support business processes and the management of the enterprise. ERP software is further used to manage company information of enterprises of various kinds in any field of technology by means of automatic data processing systems such as computers or computer systems. During the use of such software a huge amount of data is usually created, which contains important business information and which has to be archived from time to time.

According to the state of the art (see Helmut Stefani, Datenarchivierung mit SAP, Galileo Press GmbH, Bonn 2002, ISBN 3-89842-212-7), archiving can be performed automatically by archiving software tools, which can be

part of the ERP software. Such tools can consist of a writing module, which stores (writes) the data objects to be archived sequentially in archive files, and a deleting module, which deletes the successfully archived data from the original data object base. The writing module can select the data objects to be archived from the data base according to specific criteria, e.g. the creation time of the data. It usually does not modify the original data objects or data base. The deleting module staggeredly reads the archive file sequentially and deletes the data objects found in the archive file from the original data base. This ensures that only such data objects are deleted from the original data base, which are readably stored in the archive file. The time for the archiving procedure as a whole depends on the amount of data and varies from a few milliseconds to several hours or days. Consequently, there is in many cases a considerable time gap between writing the data into the archive file and deleting the data from the original data base. This time gap can be a reason for the following problems:

As long as the data objects are still available in the original data base, they can still be modified during said time gap. Because the deleting program does not compare the archived data object and the data object to be deleted, such modifications can be lost. This has not only the consequence of the loss of the amended data, it can additionally have the consequence that certain business processes can not be completed.

An other problem arises, if several archiving processes run in parallel. Then it can happen, that one data object is archived several times, and is no longer

unambiguously identifiable. This can have the consequence that evaluations or statistical analysis, which use the archive files, produce wrong results.

It can also happen that data objects in the original data base are read by the writing module and are simultaneously modified by an other software application. In such a case, the data can be transferred from an archiveable status to a non archiveable status. In consequence, data objects which are not archiveable are written into the archive file and are deleted from the original data base. In effect, this can result in a loss of data.

Thus, there is a need for a method and/or data processing system providing a more efficient solution of the problems described above.

Summary of the Invention

In accordance with the invention, as embodied and broadly described herein, methods and systems consistent with the principles of the invention provide for moving data objects in a computer system from a first to a second storage location, comprising:
a) selecting one or more data objects having an identifier (ID) from the first storage location,
b) storing said ID in a first lock object,
c) storing said ID in a second lock object,
d) storing a data object, the ID of which is contained in the first lock object, at the second storage location,
e) deleting a data object, the ID of which is contained in the first lock object, from said first storage location,

f) deleting an ID from the first lock object earliest at a time at which step c) for the respective data object assigned to that ID has been completed,

g) deleting an ID from the second lock object earliest at a time at which step b) for a particular ID has been completed.

By using this method, software applications, which require access to data objects, can check by querying the lock object, whether the data object to be accessed are subject to a moving process or not. If yes, the access to that data can be postponed until the moving is completed.

In accordance with another aspect, the invention, as embodied and broadly described herein, methods and systems consistent with the principles of the invention provide a computer system for processing data by means of or in a software application, comprising:

- memory for storing program instructions;
- input means for entering data;
- storage means for storing data;
- a processor responsive to program instructions
- programm instructions to carry out a method as of any of claims 1 to 12.

The invention and its embodiments are further directed to a computer readable medium and a carrier signal comprising instructions for processing data according to inventive method and in its embodiments.

An advantage of the invention and its embodiments is that the security against data loss in data moving and archiving procedures is greatly improved. This avoids

in consequence a lot of time and money for data retrieving.

Additional objects and advantages of the invention and its enbodiments will be set forth in part in the description, or can be learned by practice of the invention. Objects and advantages will be realized and attained by means of the elements and combinations particularly pointed out in the appended claims. Embodiments of the invention are disclosed in the detailed description section and in the dependent and appended claims as well.

It is understood that both the foregoing general description and the following detailed description are exemplary and explanatory only and are not restrictive of the invention and its embodiments, as claimed.

Brief Description of the Drawings

The accompanying drawings, which are incorporated in and constitute a part of this specification, illustrate examples of embodiments of the invention and, together with the description, explain the principles of the invention. In the drawings,

Fig. 1 is a schematic block diagram of the implementation of the inventive method within a computer system.

Fig. 2 is a schematic diagram of an exemplary structure of a data object in accordance with the principles of the inventive method.

Fig. 3 is an exemplary flow diagram of an implementation of the selecting module shown in Fig. 1.

Fig. 4 is an exemplary flow diagram of an implementation of the writing module shown in Fig. 1.

Fig. 5 is an exemplary flow diagram of an implementation of the deleting module shown in Fig. 1.

Fig. 6 is an exemplary flow chart of a further implementation of the selection and writing module mentioned in Fig. 1.

Fig. 7 shows an exemplary flow chart of how any software application can use the concept of the P- and T-locks.

Fig. 8 shows a process alternative to that shown in Fig. 7, including a conditional deletion of a P-lock.

Fig. 9 shows an example of a flow chart for a software module by means of which the locks can be deleted.

Detailed description

Computer system and program are closely related. As used hereinafter, phrases, such as "the computer provides" and "the program provides or performs specific actions", "a user performs a specific action" are convenient abbreviation to express actions by a computer system that is controlled by a program or to express that the program or program module is designed to enable the computer system to perform the specific

action or the enable a user to perform the specific action by means of a computer system.

Reference will now be made in detail to the principles of the invention by explaining the invention on the basis of the archiving process, examples of which are illustrated in the accompanying drawings. Examples, mentioned therein, are intended to explain the invention and not to limit the invention in any kind.

Within the concept of this description, the terms used shall have their usual meaning in the context of the field of data processing unless defined otherwise in the following section:

A computer system can be a stand alone computer such as a PC or a laptop or a series of computers connected as a network, e.g. a network within a company, or a series of computers connected via the internet. A data object to be archived can be any kind or type of data, e.g. numerical or textual data, image data, meta data, irrespective whether the data are implemented as whole files or parts of files or fields in tables, irrespective whether they are stored in volatile memory or nonvolatile memory. As an example, data objects according to the present invention can be implemented as one or more fields of one or more tables, particularly of tables of a relational data base system, or as objects in an object orientated programming language.

The term ERP software shall be considerer to comprise any software application that supports the business processes of an enterprise.

A storage location is volatile or nonvolatile storage means accessible by the computer system. It can be any kind of computer storage means known to one of ordinary skill, e.g. RAM, magnetical or optical storage, such as floppy disk, hard disk, MO-Disk, CD-ROM, CD RW, DVD ROM, DVD RW, etc. The first and second storage location can be identical. In this case, the archived data objects have to be stored at a place different to the place of the original data objects to be archived. The second storage location can also be implemented as a file, located anywhere in the accessible nonvolatile storage means. Such file is subsequently referred to as archive file.

An identifier (ID) is a type of data, which allows an unambiguous identification of the data object to be archived, it can be implemented for example as a number or a combination of alphanumerical characters or as a characteristic part of the data object to be archived. It is clear from that definition that a data object can have a wide variety of IDs. A lock object is a data object, in which the identifiers are stored. It can be implemented e.g. as a file on a storage means or as a data array in computer memory. The first lock object is stored advantageously in a nonvolatile storage means. The second lock object can be stored in volatile and/or nonvolatile storage means.

Fig. 1 depicts one example of an implementation of a first embodiment of the invention. Fig. 1 shows a computer system 101 comprising a computer 103 having a CPU 105, a working storage 112, in which an ERP software 111 is stored for being processed by CPU 105. The second lock object is stored in working storage 112 as well. ERP software 111 comprises program modules

106, 109, 110 for carrying out the inventive data archiving process. Computer System 101 further comprises input means 113, output means 112 for interaction with a user, and general input/output means 104, including a net connection 114, for sending and receiving data. A plurality of computer systems 101 can be connected via the net connection 114 in the form of a network 113. In this case the network computers 113 can be used as further input/output means, including the use as further storage locations. Computer system 103 further comprises a first storage means 107, in which data to be archived and the first lock object are stored, and a second storage means 108, in which the archived data are stored.

In case the program modules 106, 109, 110 are processed by CPU 105 in order to carry out the inventive process, one or more data objects stored in the first storage means 107 are selected by selection module 110. Selection module 110 stores the ID of the selected data object in the first lock object at the storage location 107 and in the second lock object in the working storage 112. Writing module 106 reads the data objects and the lock object and stores such data objects, the ID of which are contained in the lock object to the second storage location 108. Deleting module 109 then reads the archived data objects in the second storage location 108 and deletes the data objects, which it could successfully read from the orignial set of data objects in the first storage location 107. After deleting a specific data object, to which an ID was assigned, that ID is deleted from the first lock object.

In an alternative embodiment, the lock object is created by the selection module and not by the writing module.

In a second implementation of the invention, a data object to be archived comprises one or more fields of one or more tables, and the ID of the respective object comprises one or more key fields of that data object. This can best be seen from Fig. 2. In this instance, various sets of data objects are created in the form of two-dimensional data arrays, i.e. two tables having columns named field A to field X and field Y, respectively, and a certain, unspecified number of lines. A field of the array or table is defined by the name of the column and the respective line. Such field can contain data to be archived. It can alternatively contain a reference to a line of a further table. For example, in table 1 field X in line 2 contains a reference to line 3 in table 2. A data object to be archived comprises fields of one line of the respective table. If one of the fields contains a reference to a line of an other table, fields of this referenced line belong to the data object, too. In the example in Fig. 2, a data object to be archived comprises the fields of line 2 in table 1 and fields of line 3 in table 2. An ID of such a data object can be implemented by the content of one or more so-called key fields, if the combination of these key fields is unique within the respective table. In the example, the fields of "field A" and "field B" can be used as key fields for table 1, whereas field A alone is key field in table 2. Within this example, the data object has the content of the fields of columns field A and B of the respective lines as ID. The ID for the data object to be archived is stored as a first type ID in the first lock object,

named persistent lock object in Fig. 2, and as a second type ID in the second lock object, named transactional lock object. The persistent lock object is implemented as a table having two columns, the first of which contains the ID. The second type ID, ID 2, can be implemented as a one-dimensional data array stored in the working memory of the computer system. However, it can be implemented as file on a nonvolatile storage means, too. The first type ID, ID 1, is deleted after the selected data object has been deleted according to step e) of the inventive process, and the second type ID, ID 2, is deleted immediately after the time as defined in step g). Alternatively, type ID 1 IDs can be deleted after all the selected data objects have been deleted according to step e). As can be seen, both ID types have identical content, the ID of the respective lines of the data to be archived. However, this is not a necessary condition. Different contents can be used for the different ID types. The persistent lock objects further contain a column by which a filename is assigned to the ID of the data object, i.e. that data object to be archived. In the example, line 1 is archived in a file named 001, lines 2 and 3 in file 002, and line 4 in file 003.

The selection of the data object can be implemented by an automatic procedure, such as a simple query, that returns all lines having a certain field that satisfies a certain condition. For example, the procedure could return all lines in which the content of a date field pre-dates or post-dates a certain deadline. Selection can also be implemented by a user to whom a selection table is presented via a graphical user interface.

A further embodiment is characterized in that in step c) the ID is stored in the second lock object immediately after performing step a) for the respective data object. Alternatively, in step c) the ID of the selected data object is stored in the second lock object shortly before the storing process according to step d) for the data object assigned to that ID is started.

A further embodiment is characterized in that in step b) the IDs of all selected data objects are stored in the first lock object before the first storing process according to step d) is started.

In a further embodiment the invention comprises h) checking before or while performing any of steps a) to c) for a data object, whether an ID for the data object has been stored in a first lock object, and if yes, skipping at least step d) for that data object.

Additionally, the invention comprises i) checking before or while performing any of steps a) to d) for a data object, whether that data object is contained in the second storage location, and if yes, skipping at least step d) for that data object.

An other embodiment is characterized by said checking is performed by querying a first lock object.

j) in case of a failure in step d) checking, whether the data object assigned to the respective ID has been completely stored in the second storage location, and in case of no, skipping at least steps e) and f) for that data object and deleting the ID from the first lock object.

The invention is now described in more detail with reference to Figs. 3 to 5, which are schematic flow diagrams of exemplary implementations of the selecting, writing and deleting modules, respectively, as shown in Fig. 1. Within the context of this description, and particularly the Fig. 3 to 9, a first type ID is called a P-lock (permanent) and a second type ID is called a T-lock (transactional). So, setting a P- or T-lock for a selected object means to store an ID of that object in a respective lock object. The term "permanent" results for the property of the P-lock of existing permanently, as long as the data object is not yet deleted from its original storage location. The term "transactional" results from the property of the T-lock of existing only as long as a specific action (e.g. checking of archiveability) is performed on a selected data object or, in other words, of being deleted shortly after the respective action has been performed.

In the flow chart of the selecting module in Fig. 3, a data object is selected in a first step 301. Subsequently, a T-lock is set on this object in a second step 302. If the T-lock was successfully set (step 303), that is, if it did not yet exist, it is checked in step 304 whether a P-lock already exists in the selected data object. If not, the next data object is selected (step 309). The setting of the T-lock (step 302) and the check (step 303) whether it is successfully set can advantageously be implemented as one "atomic" step. This means that both steps can be executed essentially at the same time or, in other words, the time gap between both steps can be essentially zero.

Both checks (steps 303 and 304) can also be implemented by querying the respective lock objects.
If a P-lock exists, the T-lock is deleted (step 308) and the next data object is selected (step 309). If no P-lock exists, it is checked in steps 305 and 306, whether the data object is archiveable. Such checking comprises a test whether the data in the data object is readable, complete, not fraught with obvious failures, etc. If the test is successful, a P-lock is set on that data object in step 307, whereby no archive file is assigned to the data object. Then the T-lock is deleted (step 308) and the next data object is selected (step 309).

In the flow chart of the writing module in Fig. 4, a data object is selected in a first step 401. Subsequently, a T-lock is set on this object in step 402. If the T-lock was successfully set (step 403), it is checked in step 404 whether a P-lock already exists in the selected data object, whereby no file must be assigned to that data object. If the condition is not fulfilled, the T-lock is deleted in step 407, and the next data object is selected in step 408. If a P-lock exists, the data object is stored in an archive file in step 405 and the archive file is assigned to the data object in step 406, e.g. by adding the file name to the lock object as shown in Fig. 2. Subsequently, the T-lock is deleted (step 407), and the next data object is selected (step 408).

In the flow chart of the deleting module in Fig. 5, a data object that has already been archived is selected (step 501). This can be implemented by checking the archive files. If a data object has been selected and successfully read from the archive file, that data

object is deleted from the original storage location (step 502), the P-lock is deleted (step 503), and the next data object is selected (step 504).

In the exemplary flow chart of a further exemplary implementation in Fig. 6, the selecting and writing module described above are combined to one module. Accordingly, a data object is selected in a first step 601. Subsequently, a T-lock is set on this object in step 602. If the T-lock was successfully set (step 603), it is checked in step 604 whether a P-lock already exists in the selected data object. If not, the next data object is selected (step 610). If a P-lock exists on that object, the T-lock is deleted (step 609) and the next data object is selected (step 610). If no P-lock exists on that object, it is checked in step 605, whether the data object is archiveable. If this check fails (step 606), the T-lock is deleted (step 609), and the next data object is selected (step 610). If the check is positive, the data object is stored (step 605) in an archive file, a P-lock is set (step 608) with the archive file assigned, the T-lock is deleted (step 609) and the next data object is selected (step 610).

Fig. 7 shows by way of an exemplary flow chart how any software application can use the concept of the P- and T-locks to ensure that the measures, the software application is going to apply on the data object, do not influence the archiving process. A software application which is programmed to have a read and/or write access to data objects, which can be subject of an archiving process as described, comprises the following steps as shown in Fig. 7. In a first step 701, the data object is selected. Then a T-lock is set

in step 702 on that object by the application. If the T-lock is successfully set (step 703), it is checked in (step 704), whether a P-lock exists on that object, otherwise the application terminates(step 707). If a P-lock exists on that object (step 704), the T-lock is deleted (step 706), and the application terminates (step 707). If no P-lock exists, i.e. the data object is not subject to an archiving process, the application can have read/write access to the data object in a working step 705. Subsequently the application deletes the T-lock (step 706) and terminates (step 707).

Fig. 8 shows a process alternative to that shown in Fig. 7, including a conditional deletion of a P-lock. In a first step 801, the data object is selected. Then a T-lock is set on that object by the application (step 802). If the T-lock is successfully set (step 803), it is checked (step 804), whether a P-lock exists on that object, otherwise the application terminates (step 809). If no P-lock exists (step 804), i.e. the data object is not subject to an archiving process, the application can have read/write access to the data object in working step 807. Subsequently, the application deletes the T-lock (step 808) and terminates (step 809). If a P-lock exists (step 804), it is checked (step 805), whether a file is assigned to it. If a file is assigned, the application deletes the T-lock (step 808) and terminates (step 809). If no file is assigned, the P-lock is deleted (step 806), and the application can have read/write access to the data object (step 807). Subsequently, the application deletes the T-lock (step 808) and terminates (step 809).

This procedure is particularly useful, in that data objects, which are not yet stored in an archive file,

can be still altered. Consequently, they can be archived only at the next archive run.

Fig. 9 shows an example of a flow chart for a software module by means of which the locks set by the modules described above can be deleted. This can be useful in cases in which no archive files are assigned to P-locks or in which P-locks have been deleted for a user. Therein, a P-lock is nothing else than a data object and can be treated in the same way as described above. In a first step 901, a P-lock is selected. Then a T-lock is set to the P-lock in step 902. If the T-lock is successfully set (step 903), it is checked in step 904, whether the P-lock has a file assigned. If the T-lock is not set successfully, the module terminates (step 907). If the selected P-lock has no file assigned (step 904), the P-lock is deleted (step 905). Then the T-lock is deleted (step 906), and the module terminates (step 907). Alternative to the termination (step 907), a next P-lock can be selected.

Modifications and adaptations of the present invention will be apparent to those skilled in the art from consideration of the specification and practice of the invention disclosed herein. The foregoing description of an implementation of the invention has been presented for purposes of illustration and description. It is not exhaustive and does not limit the invention to the precise form disclosed. Modifications and variations are possible in light of the above teachings or can be acquired from the practicing of the invention. For example, the described implementation includes software, but systems and methods consistent with the present invention can be implemented as a combination of hardware and software or in hardware

alone. Additionally, although aspects of the present invention are described for being stored in memory, one skilled in the art will appreciate that these aspects can also be stored on other types of computer-readable media, such as secondary storage devices, for example, hard disks, floppy disks, or CD-ROM; the Internet or other propagation medium; or other forms of RAM or ROM. It is intended that the specification and examples be considered as exemplary only, with a true scope and spirit of the invention being indicated by the following claims.

Computer programs based on the written description and flow charts of this invention are within the skill of an experienced developer. The various programs or program modules can be created using any of the techniques known to one skilled in the art or can be designed in connection with existing software. For example, programs or program modules can be designed in or by means of ® Java, C++, HTML, XML, or HTML with included Java applets or in SAP R/3 or ABAP.

What is claimed is:

1. A method for moving data objects in a computer system from a first to a second storage location, comprising:
a) selecting one or more data objects having an identifier (ID) from the first storage location,
b) storing said ID in a first lock object,
c) storing said ID in a second lock object,
d) storing a data object, the ID of which is contained in the first lock object, at the second storage location,
e) deleting a data object, the ID of which is contained in the first lock object, from said first storage location,
f) deleting an ID from the first lock object earliest at a time at which step c) for the respective data object assigned to that ID has been completed,
g) deleting an ID from the second lock object earliest at a time at which step b) for a particular ID has been completed.

2. The method of claim 1; wherein
a data object comprises one ore more fields of one or more tables and wherein the ID comprises one or more key fields of the one or more tables.

3. The method of claim 1 or 2, wherein
in step d) the data objects are stored in one or more files and wherein an assignment of the ID to the file or to a name of the file, in which the data object assigned to said ID is stored, is stored in the first lock object.

4. The method of one of claims 1 to 3, wherein the first lock object is stored on a nonvolatile storage means.

5. The method of one of claims 1 to 4, wherein in step c) the ID is stored in the second lock object immediately after performing step a) for the respective data object.

6. The method of one of claims 1 to 4, wherein in step c) the ID of the selected data object is stored in the second lock object shortly before the storing process according to step d) for the data object assigned to that ID is started.

7. The method of one of claims 1 to 6, wherein in step b) the IDs of all selected data objects are stored in the first lock object before the first storing process according to step d) is started.

8. The method one of claims 1 to 7, further comprising:
h) checking before or while performing any of steps a) to c) for a data object, whether an ID for the data object has been stored in a first lock object, and if yes, skipping at least step d) for that data object.

9. The method of one of claims 1 to 8, further comprising:
i) checking before or while performing any of steps a) to d) for a data object, whether that data object is contained in the second storage location, and if yes, skipping at least step d) for that data object.

10. The method of claim 9, wherein said checking is performed by querying a first lock object.

11. The method of one of claims 1 to 10, further comprising:
j) in case of a failure in step d) checking, whether the data object assigned to the respective ID has been completely stored in the second storage location, and in case of no, skipping at least steps e) and f) for that data object and deleting the ID from the first lock object.

12. The method of one of claims 1 to 11 for use in an enterprise resource planning software.

13. A computer system for processing data by means of or in a software application, comprising:
- memory for storing program instructions;
- input means for entering data;
- storage means for storing data;
- a processor responsive to program instructions
- programm instructions to carry out a method as of any of claims 1 to 12.

14. A computer program comprising program code means for performing a method as of any of claims 1 to 12 if said program is executed on a computer system.

15. A computer readable medium comprising program code for performing a method as of any of claims 1 to 12 if said program code is executed on a computer system.

16. A computer program product comprising a computer readable medium according to claim 15.

17. A computer data signal embodied in a carrier wave comprising:
program code for performing a method as of any of claims 1 to 12 if said program code is executed on a computer system.

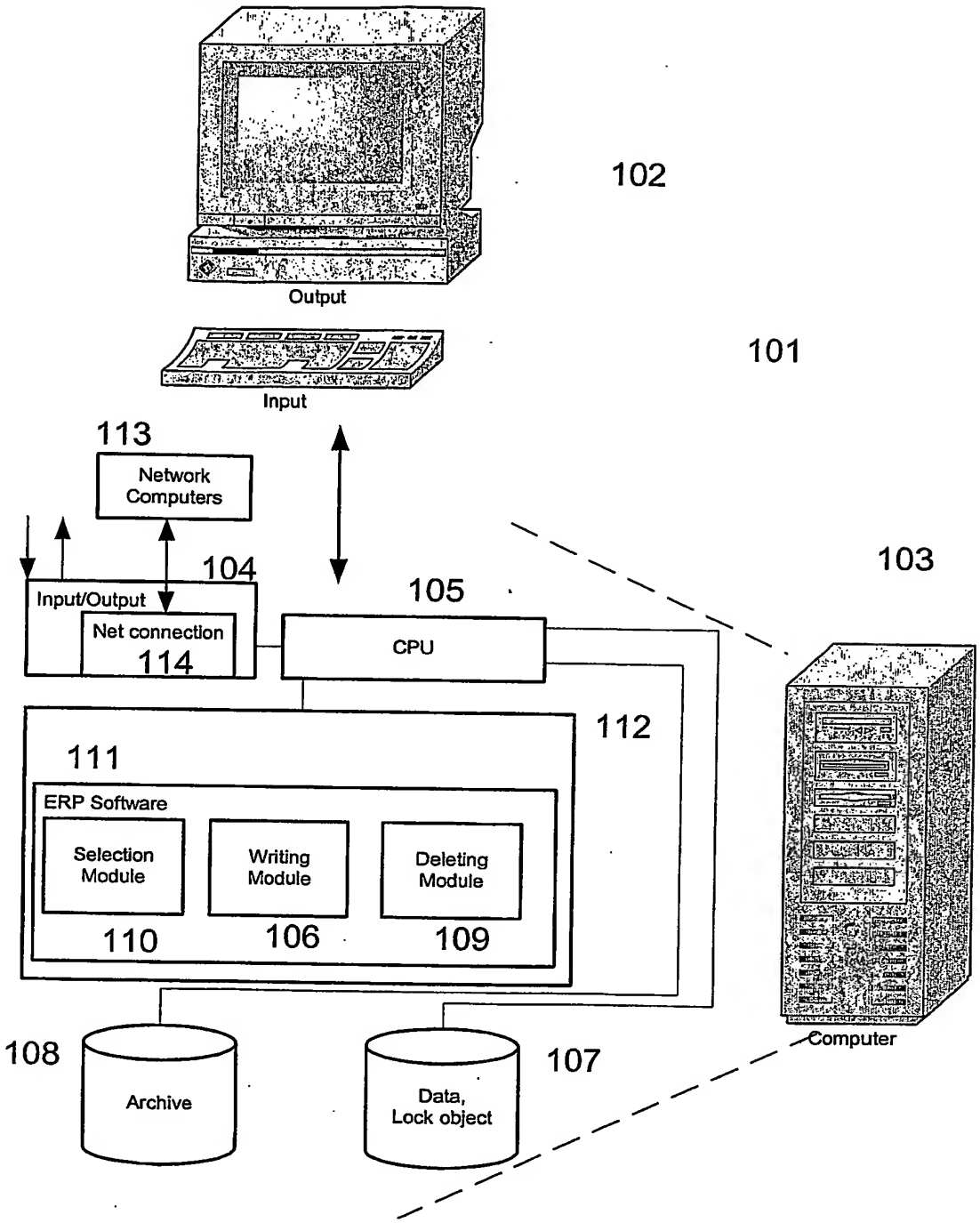
102
Output
101
Input
113
Network Computers
104
Input/Output
Net connection
114
105
CPU
103
112
111
ERP Software
Selection Module
Writing Module
Deleting Module
110
106
109
108
Archive
107
Data, Lock object
Computer

Fig. 1

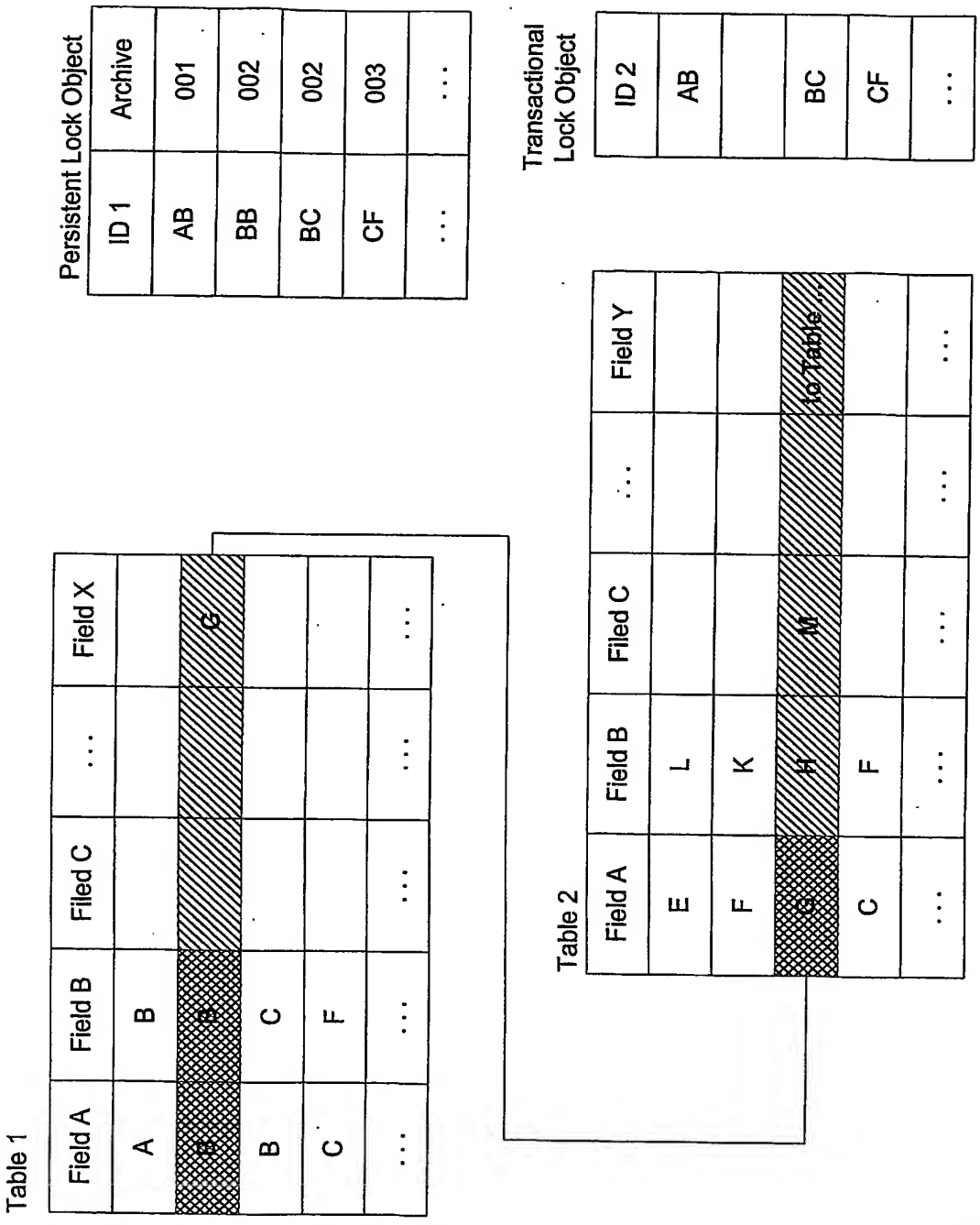
Persistent Lock Object
ID 1
Archive
AB
001
BB
002
BC
002
CF
003
Transactional Lock Object
ID 2
AB
BC
CF
Table 1
Field A
Field B
Filed C
Field X
A
B
B
C
B
C
F
C
Table 2
Field A
Field B
Filed C
Field Y
E
L
F
K
G
H
C
F
M
to Table ...

Fig. 2

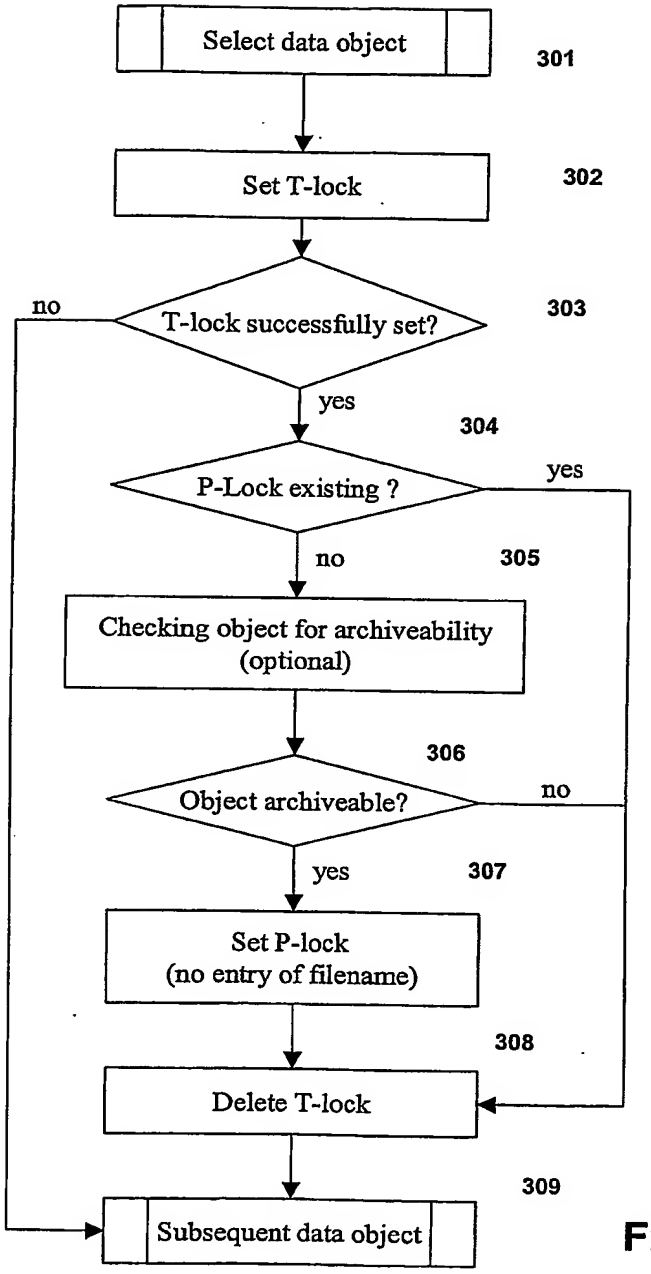
Select data object
301
Set T-lock
302
T-lock successfully set?
303
no
yes
304
P-Lock existing ?
yes
no
305
Checking object for archiveability
(optional)
306
Object archiveable?
no
yes
307
Set P-lock
(no entry of filename)
308
Delete T-lock
309
Subsequent data object

Fig. 3

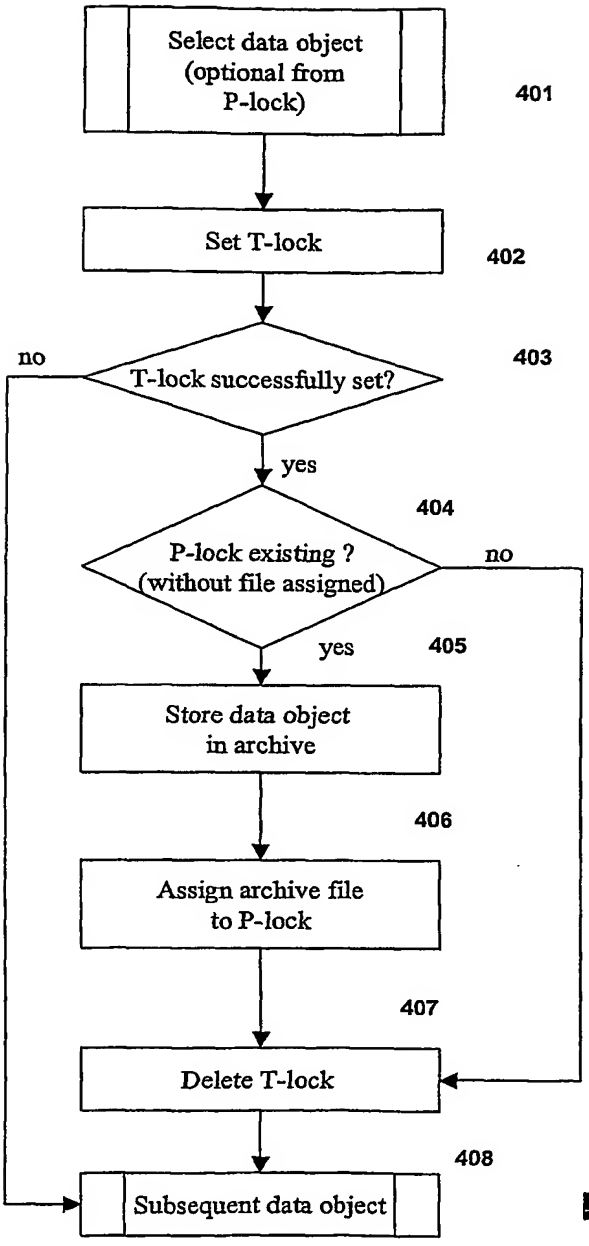
Select data object (optional from P-lock)
401
Set T-lock
402
T-lock successfully set?
403
no
yes
P-lock existing ? (without file assigned)
404
no
yes
405
Store data object in archive
406
Assign archive file to P-lock
407
Delete T-lock
408
Subsequent data object

Fig. 4

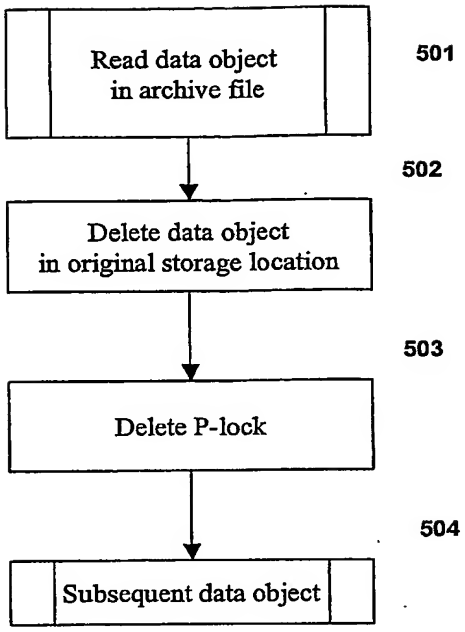
Read data object in archive file
501
502
Delete data object in original storage location
503
Delete P-lock
504
Subsequent data object

Fig. 5

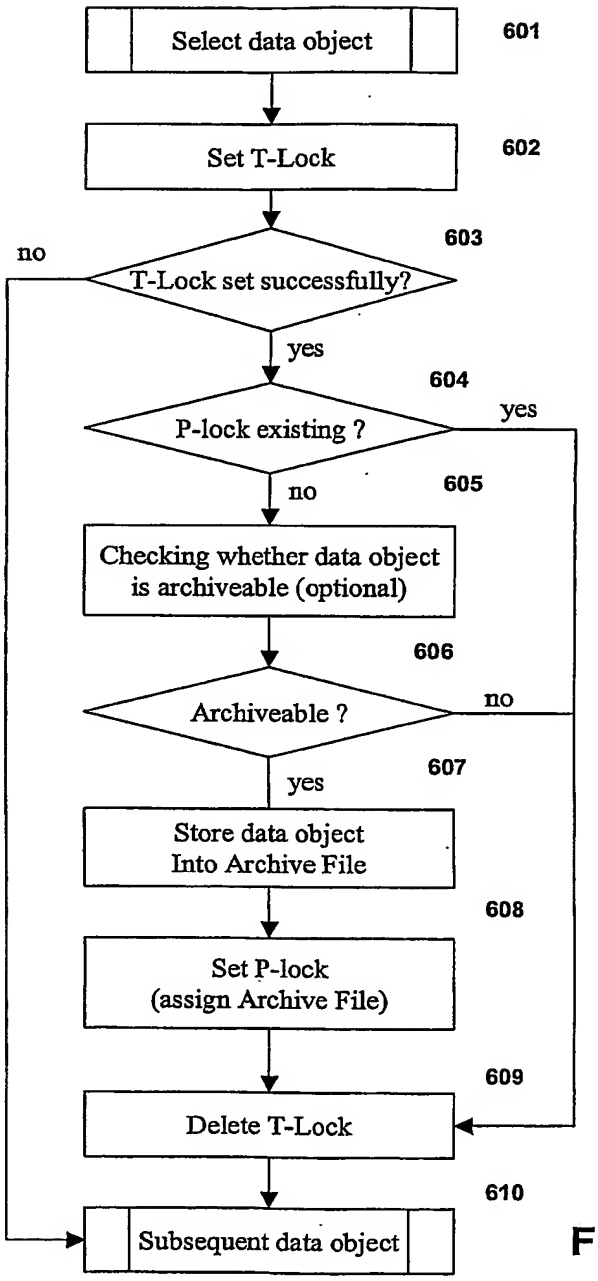
Select data object
601
Set T-Lock
602
603
no
T-Lock set successfully?
yes
604
yes
P-lock existing ?
no
605
Checking whether data object is archiveable (optional)
606
Archiveable ?
no
yes
607
Store data object Into Archive File
608
Set P-lock (assign Archive File)
609
Delete T-Lock
610
Subsequent data object

Fig. 6

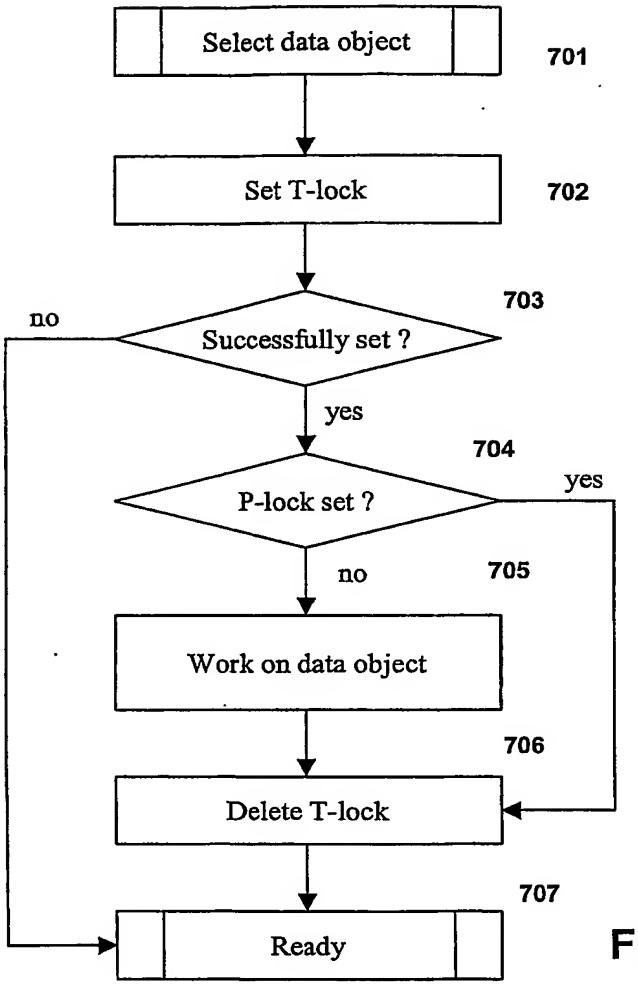
Select data object
701
Set T-lock
702
703
no
Successfully set ?
yes
704
yes
P-lock set ?
no
705
Work on data object
706
Delete T-lock
707
Ready

Fig. 7

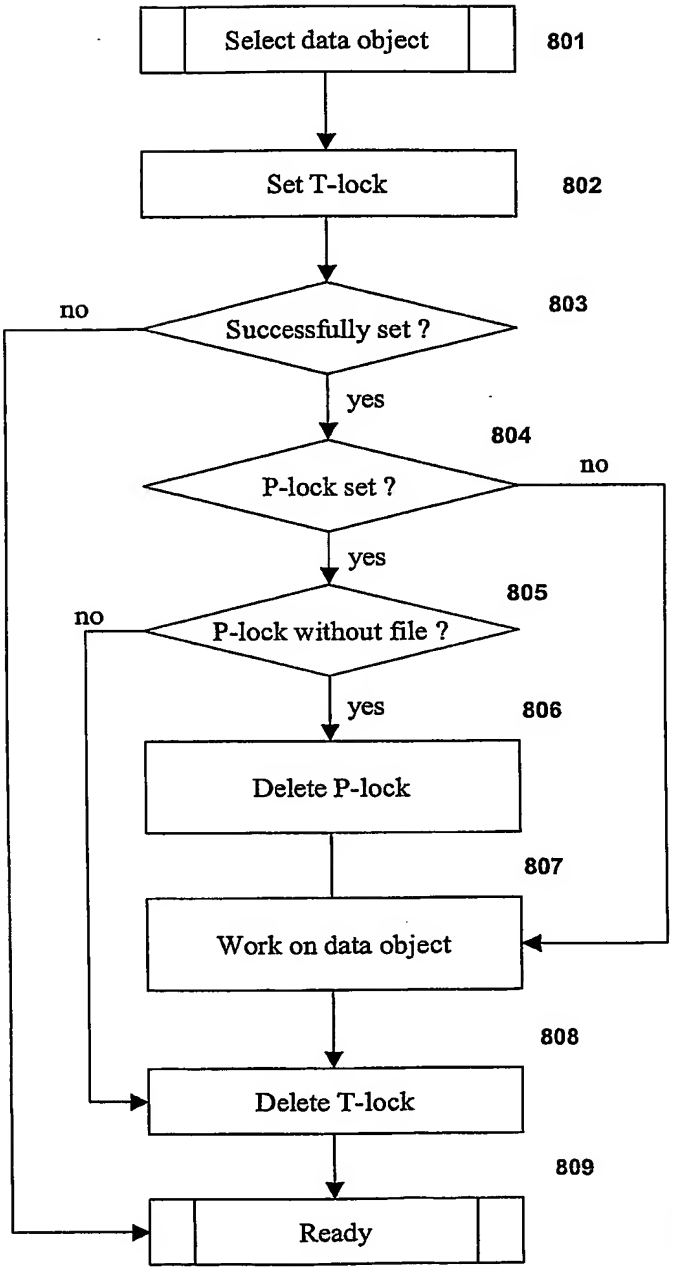
Select data object
801
Set T-lock
802
Successfully set ?
803
no
yes
P-lock set ?
804
no
yes
P-lock without file ?
805
no
yes
806
Delete P-lock
807
Work on data object
808
Delete T-lock
809
Ready

Fig. 8

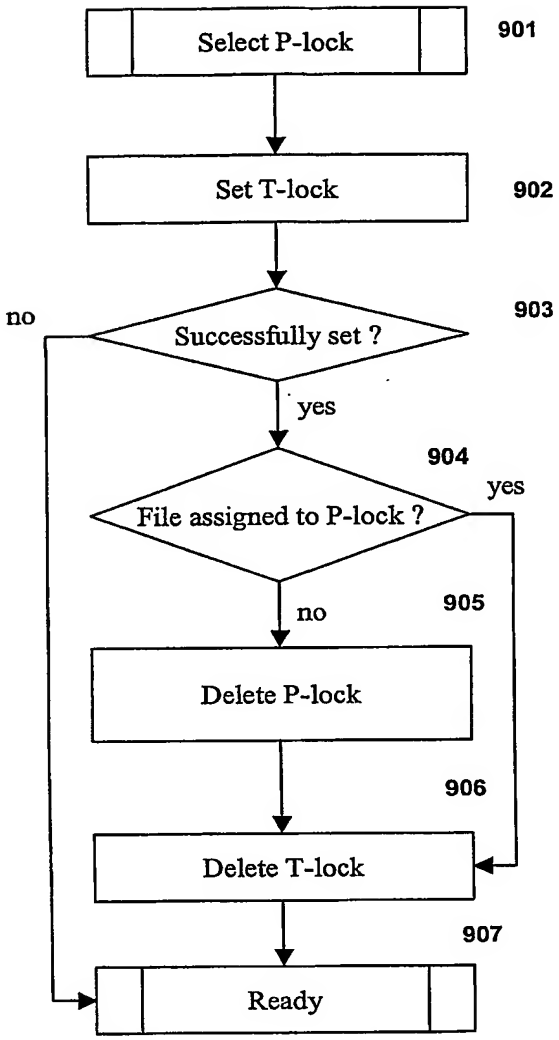
Select P-lock
901
Set T-lock
902
Successfully set ?
903
no
yes
File assigned to P-lock ?
904
yes
no
905
Delete P-lock
906
Delete T-lock
907
Ready

Fig. 9

INTERNATIONAL SEARCH REPORT

International Application No
PCT/EP 03/09832

A. CLASSIFICATION OF SUBJECT MATTER
IPC 7 G06F17/30 G06F9/46

According to International Patent Classification (IPC) or to both national classification and IPC

B. FIELDS SEARCHED

Minimum documentation searched (classification system followed by classification symbols)
IPC 7 G06F

Documentation searched other than minimum documentation to the extent that such documents are included in the fields searched

Electronic data base consulted during the international search (name of data base and, where practical, search terms used)
EPO-Internal

C. DOCUMENTS CONSIDERED TO BE RELEVANT

Category °	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
X	STEFANI H.: "Datenarchivierung mit SAP" May 2002 (2002-05) , SAP PRESS, GALILEO PRESS , BONN XP002266517 ISBN: 3-89842-212-7 cited in the application page 35 -page 40 page 57 page 63 -page 75 page 84 -page 85 page 211 -page 212 ---	1-17
A	US 5 566 319 A (LENZ NORBERT) 15 October 1996 (1996-10-15) column 1, line 66 -column 2, line 23 --- -/--	1-17

[X] Further documents are listed in the continuation of box C.

[X] Patent family members are listed in annex.

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"A" document defining the general state of the art which is not considered to be of particular relevance
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"O" document referring to an oral disclosure, use, exhibition or other means
"P" document published prior to the international filing date but later than the priority date claimed

"T" later document published after the international filing date or priority date and not in conflict with the application but cited to understand the principle or theory underlying the invention
"X" document of particular relevance; the claimed invention cannot be considered novel or cannot be considered to involve an inventive step when the document is taken alone
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"&" document member of the same patent family

Date of the actual completion of the international search	Date of mailing of the international search report
12 January 2004	27/01/2004

Name and mailing address of the ISA	Authorized officer
European Patent Office, P.B. 5818 Patentlaan 2 NL – 2280 HV Rijswijk Tel. (+31–70) 340–2040, Tx. 31 651 epo nl, Fax: (+31–70) 340–3016	Huber, A

Form PCT/ISA/210 (second sheet) (July 1992)

C.(Continuation) DOCUMENTS CONSIDERED TO BE RELEVANT

Category °	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
A	ANDREW S. TANENBAUM: "Modern Operating Systems" 1992 , PRENTICE-HALL INTERNATIONAL, INC. , NEW JERSEY, USA XP002266518 ISBN: 0-13-595752-4 page 494 -page 496 ---	1-17
P,A	EP 1 283 477 A (SAP AG) 12 February 2003 (2003-02-12) the whole document ---	1-17
A	EP 0 499 422 A (IBM) 19 August 1992 (1992-08-19) page 3, line 1 - line 33 -----	1-17

INTERNATIONAL SEARCH REPORT

Information on patent family members

International Application No
PCT/EP 03/09832

Patent document cited in search report	Publication date	Patent family member(s)	Publication date
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		JP 6187232 A	08-07-1994
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		WO 03014968 A1	20-02-2003
EP 0499422 A	19-08-1992	EP 0499422 A2	19-08-1992
		JP 4310148 A	02-11-1992
		JP 8027755 B	21-03-1996
		US 5327556 A	05-07-1994